

ACM ICPC Manila 2016

Solution Sketches

Problem G: Go Go Go Special Action Force!

- Easy problem
- Check if the input is a valid Sudoku grid
 - Option 1: sets
 - Option 2: bitmasks
- Just be careful!



Problem L: LoL Tournament

- The tournament is a binary tree
- The highest chance is those with fewest battles
 - i.e., minimum depth
- Construct tree and compute depths
- Special case: $p = 0$ or $p = 1$
 - All nodes are equiprobable!




Problem F: Frog Pushers

- Given sequence of edge removals, compute **MST** before *each* removal
- *Bad:*
 - E times Prim's/Kruskal's
 - Too slow
 - **$O(E^2 \log E)$**



Problem F: Frog Pushers

- *Good:*
 - Reverse the input; so we're now dealing with edge additions.
 - After each edge addition, compute new MST. Throw away unused edges *permanently*.
 - This keeps edge count at most **$V-1$** always
 - **$O(EV \log V)$**
- 

Problem F: Frog Pushers

- *Better:*
 - Maintain cost-sorted list of edges.
 - **$O(EV \alpha(V))$**
- *Betterer:*
 - MST update after an insertion can be done in **$O(V)$**
 - BFS to find the largest cost in the (new) cycle
 - **$O(EV)$**

Problem F: Frog Pushers

- Asymptotically faster solutions exist
 - (I think)
 - Link-cut trees?



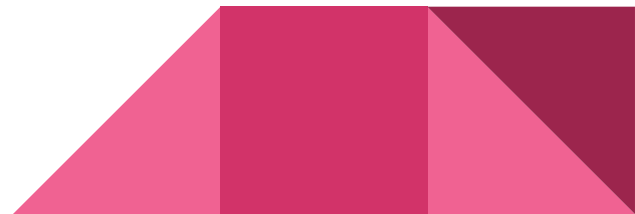
Problem K: Off the Rails

- Assume single segment / line.
 - Calculus optimization / linear algebra
- Let $f(i, j)$ be this answer for the segment $pt[i..j]$.
- $f(i, j)$ can be computed in $O(n)$.



Problem K: Off the Rails

- Assume single segment / line.
 - Calculus optimization / linear algebra
- Let $f(i, j)$ be this answer for the segment $pt[i..j]$.
- $f(i, j)$ can be computed in **$O(1)$** by **precomputing some sums.**



Problem K: Off the Rails

- Let $F(j)$ be the answer for $pt[1..j]$.
 - $F(0) = 0$
 - $F(j) = \min_i F(i) + C + f(i+1, j) [0 \leq i < j]$
- Dynamic programming in **$O(n^2)$**



Problem K: Off the Rails

- Yes, the title doesn't start with a "K"... it's **off the rails**



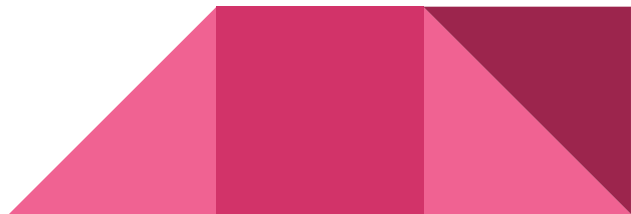
Problem B: Balloon Distribution

- Simple solution: Priority queue
 - **$O(N \log M)$** . Too slow



Problem B: Balloon Distribution

- Assume $\mathbf{a/b}$ is the ratio in the last rank. How many balloons?
 - $\text{count}(a/b) := \sum_{1 \leq i \leq M} \text{floor}(P[i]*b/a)$
- **Binary search** to find the real last ratio $\mathbf{a/b}$.
- Quite tricky. Be careful about precision.

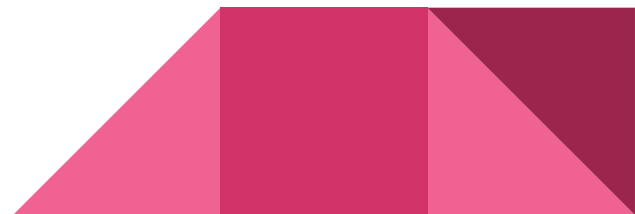


Problem B: Balloon Distribution

- One way to do it:
 - Let $M = \max(P[i]) + 1$. Find smallest D : $\text{count}(M/D) \geq N$.
 - *Binary search*
 - Next, let $d[i] = \text{floor}(P[i] * D / M)$
 - Sort the list $[P[1]/d[1], P[2]/d[2], \dots, P[M]/d[M]]$
 - The answer is in this list. *Binary search* again
 - No big ints/floats necessary. Use 64-bits.
- **$O(M (\log M + \log N))$**

Problem J: Jack and Jill and Joe

- The dimensions are a^2 and $b(b+a)$, and we want
 - $(a^2 - b(b+a))^2 \leq 900$
 - equivalently, $a^2 - b(b+a) = d$ for some $|d| \leq 30$.
 - equivalently, $(2b+a)^2 - 5a^2 = 4d$ for some $|d| \leq 30$.
- Gives us 61 generalized Pell equations!



Problem J: Jack and Jill and Joe

- Solutions follow a familiar pattern: a 2D recurrence.
 - In fact, Fibonacci recurrence:
 - (a, b) is a solution $\rightarrow (a + b, a)$ is a solution
- Pattern-match and precompute all solutions for all d .



Problem J: Jack and Jill and Joe

- There are so few solutions: 443
 - they grow exponentially
- To answer a query **N**, easiest to just linear search.
- Watch out for d's with multiple “solution families”!



Problem I: Imelda's Shopping Spree

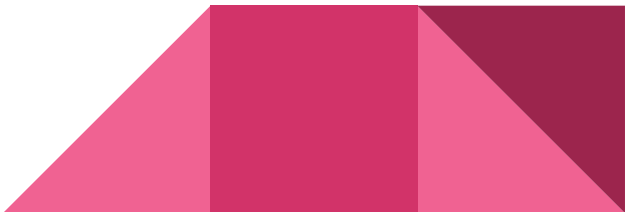
- Range queries:
 - Range increment.
 - Range reverse.
 - Query: How many contiguous increasing subsequences?




Problem I: Imelda's Shopping Spree

- Compute the difference array $D[i] = P[i+1] - P[i]$.

Then:

- Range increment reduces to 2 “point increments”.
 - Range reverse reduces to range reverse + range negate + 2 “point increments”.
 - Query reduces to counting the number of contiguous positive subsequences.
- 

Problem I: Imelda's Shopping Spree

- **Splay tree** or **treap** can handle reverse. Need the following info for every node:
 - Sum of values.
 - Longest pos./neg. streak on left and right.
 - Number of contiguous pos./neg. subsequences.
 - Need to keep track of negatives to account for *range negates*
 - “Reverse” flag for lazy propagation
- 

Problem C: Convex Quadrilaterals

- Only **convex hull** matters.
- **DP**: Choose starting edge and keep track of current edge and number of edges used so far.
- **$O(N^3)$** (**$O(N^2)$** states, **$O(N)$** transition)
- To compute (signed) areas, you can use *shoelace formula*



Problem E: Expression

- Easy/tricky cases
 - If $y = 0$, output $xx-$
 - If $y > 0$ and $x = 0$, impossible.
- General case: 28 ops quite tight!
 - Binary, ternary, etc., don't work



Problem E: Expression

- Let $\{y\}$ be a string that evaluates to xy for any $x > 0$.
 - Output “ $\{y\}_x/$ ”
- You can write $\{y\}$ in ≤ 27 operations!
 - for $0 < y < 1212$



Problem E: Expression

- Tricks:
 - $\{1\} = \times$
 - $\{ab\} \leq \{a\}\{b\}_{\times/*}$ for $a, b > 1$
 - $\{a+b\} \leq \{a\}\{b\}_+$ for $a, b > 0$
- “ \leq ” means “has the ff. candidate solution”



Problem E: Expression

- Tricky trick:
 - Use subtraction!
 - $\{ab-c\} \leq \{a\}\{b\} \times / * \{c\} -$
 - for $a, b, c < ab-c$.
 - Without it, you might fail 1103.
- With these tricks, generate all.



Problem E: Expression

- Worst case?
 - **27** operations: 823, 1006, 1111, 1198, 1211 ...
 - **28** operations: 1114, 1138, 1166, 1193
- I'd like to see better solutions!



Problem E: Expression

- Be careful with intermediate values!
 - They should be small
 - Always divide first:
 - $\{a\} \{b\} * x / \rightarrow \{a\} \{b\} x / *$



Problem A: AdoraBalls

- Invert the 4D matrix, and check if the resulting product with the vector is all nonnegative.
- Tricky when non-invertible.



Problem A: AdoraBalls

- Another way to look at it:
 - Given 4D points P_0, P_1, P_2, P_3, P_4 , can P_0 be expressed as a *nonnegative* rational linear combination of P_1, P_2, P_3, P_4 ?
 - Also known as *convex combination*



Problem A: AdoraBalls

- Project all points to a 3D-hyperplane
 - $x+y+z+w = 1$
- The problem is now: Is the 3D point P_0' inside the *tetrahedron* with vertices P_1', P_2', P_3', P_4' ?



Problem A: AdoraBalls

- Is the 3D point P_0' inside the *tetrahedron* with vertices P_1', P_2', P_3', P_4' ?
- General case:
 - For each vertex P_i' , check if P_0' is in the same side of the plane determined by the remaining three vertices.
 - Cross products + Dot products

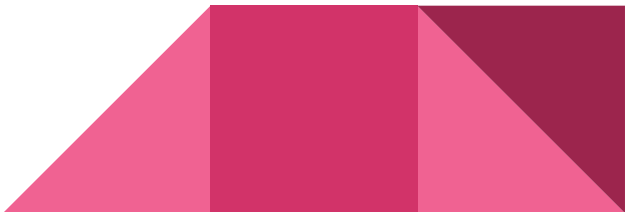


Problem A: AdoraBalls

- Watch out for special/degenerate cases! E.g.
 - The tetrahedron has zero area, e.g. it's 2D, 1D, 0D, etc.
 - One (or more) of the original 4D points is the origin $(0,0,0,0)$.
 - Can't be projected; needs special handling.
- Be careful with precision!
 - Add fractions with lcm.



Problem D: Disco Dance Debacle

- A disco dance always alternates row and column and ends up at the starting point.
 - Thus, number of lit cells it unlights per row/column is even.
 - Thus, *necessary* condition is: each row/column must have even no. of lit cells.
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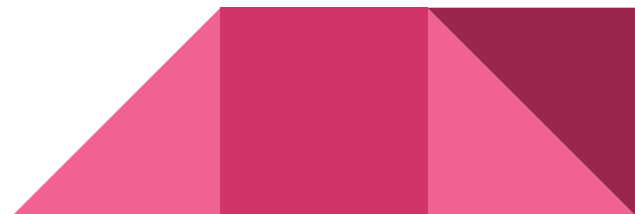
Problem D: Disco Dance Debacle

- It's also *sufficient*!
 - Draw a bipartite graph with $m+n$ nodes, and edge (i, j) if cell (i, j) is lit.
 - Then a disco dance is just a simple cycle!
 - Each connected component can be decomposed into a cycle
 - **Eulerian cycle**



Problem D: Disco Dance Debacle

- So the goal is: make the no. of lit cells even, per row/column.
- If there are **R** rows and **C** columns with odd no. of lit cells, then **max(R,C)** moves are *necessary* to make them all even.
- It's also *sufficient*!
 - Why?



Problem D: Disco Dance Debacle

- So we have to compute **R** and **C**. (Symmetric)
- Sweep line + Range queries to process events:
 - Range increment/decrement.
 - Range “how many are equal to 0”.
- **Segment tree:**
 - Insight: “0” only ever appears as the min.
 - (min value, frequency of min value)



Problem H: Handbags

- BFS/DFS too slow; large dimensions
- Be careful; prices can't "pass through" sources.
- Sometimes sources even block off some sections of the map.



Problem H: Handbags

- **Solution: Coordinate compression.**
 - Let $X = \{0, a\} \cup \{x-1, x, x+1 \text{ for all sources } (x,y)\}$
 - Let $Y = \{0, b\} \cup \{y-1, y, y+1 \text{ for all sources } (x,y)\}$
 - Dijkstra on the coordinates X and Y .
 - Coordinate compression
 - “Fill in” each (big) cell in $O(1)$ using tricky arithmetic



Thank you!

- Credits

- **AdoraBalls** - Atienza
- **Balloon Distribution** - Muga, Atienza
- **Convex Quadrilateral** - Manalastas
- **Disco Dance Debacle** - Atienza
- **Expression** - Atienza
- **Frog Pushers** - Atienza
- **Go Go Go Special Action Force!** - Chua
- **Handbags** - Muga, Atienza
- **Imelda's Shopping Spree** - Atienza
- **Jack and Jill and Joe** - Sioson, Atienza
- **Off the Rails** - Zuniga
- **LoL Tournament** - Atienza

- Judges

- Dr. Allan Sioson
 - Also chief judge
- Dr. Pablo Manalastas
- Mr. Kevin Atienza
 - Also testing
 - Also additional test cases
 - Also problem extensions
- Dr. Philip Zuniga
- Dr. Felix Muga
- Dr. Caslon Chua